Computing inferences for large-scale continuous-time Markov chains by combining lumping with imprecision

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Yet another talk about imprecise Markov chains?

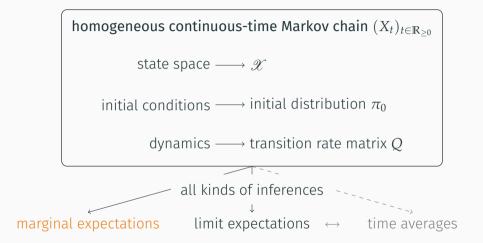


Yet another talk about imprecise Markov chains?

TeaserAllowing for more imprecision results in faster computations!



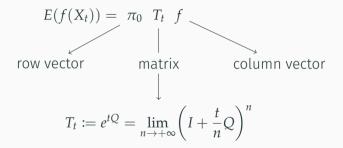
Continuous-time Markov chain refresher



Marginal expectation

Fix some $f: \mathscr{X} \to \mathbb{R}$ and $t \in \mathbb{R}_{\geq 0}$.

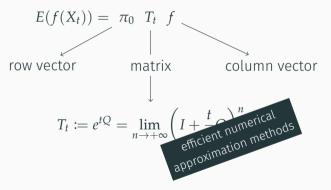
It is well-known that



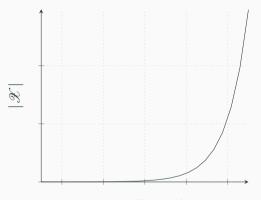
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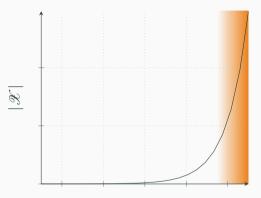


State space explosion



system dimensions

State space explosion



system dimensions

problem computing $E(f(X_t))$ becomes intractable

solution reduce the number of states *somehow*

Lumping

Informally

Taking together (lumping, aggregating) states yields the lumped stochastic process, which has a (significantly) smaller state space.

Formally

The lumped state space $\hat{\mathscr{X}}$ is a partition of \mathscr{X} . $[1 < |\hat{\mathscr{X}}| \ll |\mathscr{X}|]$

The lumping map $\Lambda \colon \mathscr{X} \to \mathscr{X}$ maps states to corresponding lumps. [A is a surjection]

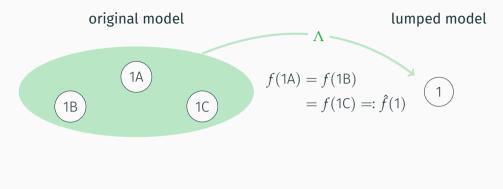
The *lumped stochastic process* is defined as

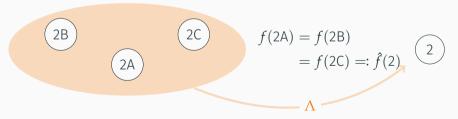
 $\hat{X}_t \coloneqq \Lambda(X_t)$ for all $t \in \mathbb{R}_{\geq 0}$.

original model







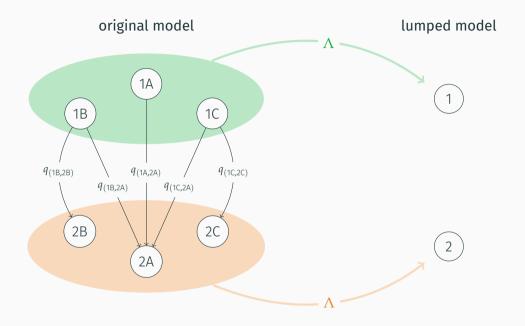


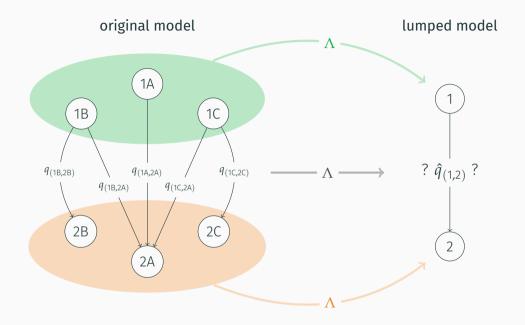
Fix some $f: \mathscr{X} \to \mathbb{R}$ and $\hat{f}: \mathscr{X} \to \mathbb{R}$ such that $\hat{f} \circ \Lambda = f$.

Then since $\hat{X}_t = \Lambda(X_t)$,

 $E(f(X_t)) = E(\hat{f}(\hat{X}_t)).$

If we cannot compute $E(f(X_t))$ tractably, can we compute $E(\hat{f}(\hat{X}_t))$ tractably?





In general, the lumped stochastic process $(\hat{X}_t)_{t\in\mathbb{R}_{\geq 0}}$

- has dynamics that cannot be determined "immediately" from Q and Λ , and
- is not homogeneous nor Markov.

In general, the lumped stochastic process $(\hat{X}_t)_{t \in \mathbb{P}}$

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- is not homogene We cannot tractably compute $E(\hat{f}(\hat{X}_t))$. In general, the lumped stochastic process $(\hat{X}_t)_{t \in \mathbb{P}}$

Lety from Q and Λ , and

We therefore consider the set $\mathbb{P}_{\pi_0,O,\Lambda}$ of consistent stochastic processes, that

- contains the lumped stochastic process.
- + takes the form of an imprecise continuous-time Markov chain and
- + is fully parameterised by

 $\hat{\pi}_0$ the lumped initial distribution, and

 $\hat{\mathscr{Q}}$ the set of *possible* lumped transition rate matrices. which in turn are fully defined by π_0 , Q and A.

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We therefore consider the set $\mathbb{P}_{\pi_0,O,\Lambda}$ of consistent stochastic processes, that

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We can tractably compute lower and upper bounds for $E(\hat{f}(\hat{X}_t))$. \sim the set of possible lumped transition rate matrices,

which in turn are fully defined by π_0 . O and A.

Essential to our computations are the lumped initial distribution $\hat{\pi}_0$, given by

$$\hat{\pi}_0(\hat{x}) = \sum_{x \in \Lambda^{-1}(\hat{x})} \pi_0(x) ext{ for all } \hat{x} \in \mathscr{X}$$
 ,

and the lower transition rate operator \hat{Q} that, for all $g \in \mathscr{L}(\hat{\mathscr{X}})$, is given by

$$[\underline{\hat{Q}g}](\hat{x}) = \min\left\{\sum_{\hat{y}\in\hat{\mathscr{X}}} g(\hat{y}) \sum_{y\in\Lambda^{-1}(\hat{y})} Q(x,y) \colon x\in\Lambda^{-1}(\hat{x})\right\} \text{ for all } \hat{x}\in\hat{\mathscr{X}}.$$

 $[\Lambda^{-1} \text{ is the set-valued inverse of } \Lambda]$

Note: in practice, this optimisation usually simplifies considerably

 $\hat{\pi}_0 \quad \underline{\hat{T}}_t \quad \hat{f}$

Fix some
$$f: \mathscr{X} \to \mathbb{R}$$
 and $\hat{f}: \mathscr{X} \to \mathbb{R}$ such that $\hat{f} \circ \Lambda = f$.
Then

$$\leq E(\hat{f}(\hat{X}_{t})) = E(f(X_{t})) = \pi_{0} T_{t} f$$

$$T_{t} = \lim_{n \to +\infty} \left(I + \frac{t}{Q} O roximation \\ numerical approximation \\ numerical approxima$$

Fix some
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Then
 $\hat{\pi}_0 \quad \underline{\hat{T}}_t \quad \hat{f} \leq E(\hat{f}(\hat{X}_t)) = E(f(X_t)) = \pi_0 \quad T_t \quad f$
row non-linear column
vector vector operator vector
 $\hat{T}_t := \lim_{n \to +\infty} \left(I + \frac{t}{n} \underline{\hat{Q}} \right)^n$
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numerical approximation
numerical approximation
methods are intractable

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efficient numerical
approximation methods
 $T_t = \lim_{n \to +\infty} \left(I + \frac{t}{t} \stackrel{\Lambda}{A} \right)^n$
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numerical approximation

Fix some $f \colon \mathscr{X} \to \mathbb{R}$ and $\hat{f} \colon \mathscr{\hat{X}} \to \mathbb{R}$ such that $\hat{f} \circ \Lambda = f$.

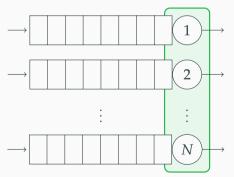
Then

$$\hat{\pi}_0 \quad \underline{\hat{T}}_t \quad \hat{f} \leq E(\hat{f}(\hat{X}_t)) = E(f(X_t)) = \pi_0 \quad T_t \quad f \leq -\hat{\pi}_0 \quad \underline{\hat{T}}_t \quad (-\hat{f}).$$
conjugacy

A queueing example » Set up

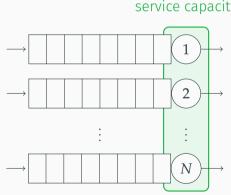
(Cardoen, 2018) studies a system of parallel queues with *processor sharing*.

shared service capacity



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shared service capacity

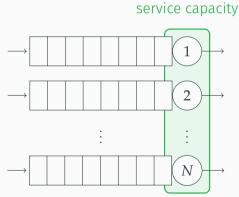
Full system description

• state is (i_1, i_2, \ldots, i_N)

$$- |\mathscr{X}| = (K+1)^N$$

A queueing example » Set up

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shared Full system description

- state is (i_1, i_2, \ldots, i_N)
- $|\mathscr{X}| = (K+1)^N$

Only interested in i_1

- (i_1, i_2, \dots, i_N) contains too much information
- ¿ lumped state space ?

A queueing example » Trade-off between imprecision and computation time

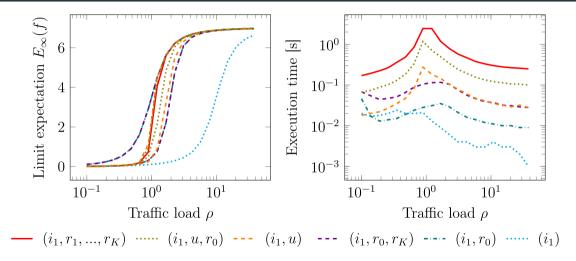


Figure taken from (Cardoen, 2018)

References

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